

Computing the Character Table of a 2-local Maximal Subgroup of the Monster

Anthony Pisani

Monash University

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Outline

Aim

Compute the character table of $2^{5+10+20} \cdot (S_3 \times L_5(2))$.

1 Background

2 Character Table Computation

Motivation

Character tables nicely “summarise” groups.

All groups factor into a composition series of simple groups, so their character tables are of particular interest.

Whole books have been published on the subject (e.g. the *ATLAS*)!

Simple groups still have a subgroup structure: look at their *maximal* subgroups too.

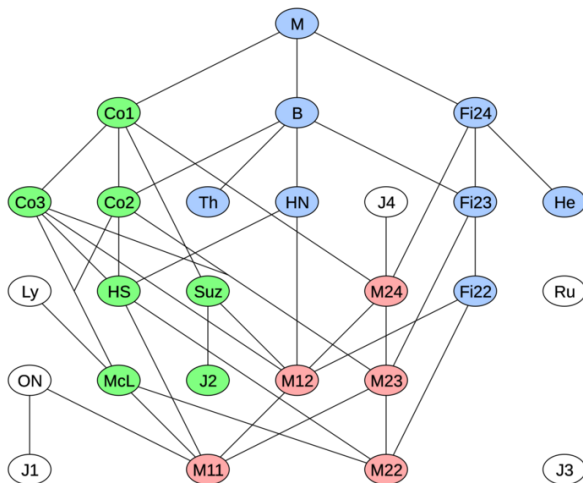
Sporadic groups & the Monster

The **CFSG** (Classification of Finite Simple Groups) asserts all simple groups, except 26 **sporadic groups**, belong to one of 18 infinite families.

The Monster \mathbb{M} is the largest sporadic group by far: $\approx 2 \times 10^{20}$ times bigger than the Baby Monster \mathbb{B} in second place. First constructed by Griess in 1982 as a 196884-dimensional matrix group.

$$|\mathbb{M}| \approx 8 \times 10^{53}.$$

Sporadic groups & the Monster



The Sporadic Groups (SporadicGroups.png, Wikimedia Commons: CC BY-SA 4.0)

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But the problems for the Monster's maximal subgroups remained: even the list of subgroups wasn't settled until 2024.

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As a permutation group: $\approx 10^{20}$ points (vs. \mathbb{B} 's $\approx 10^{10}$)

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Bray and Wilson (2006) give permutation representations for all then-known maximal subgroups of \mathbb{M} , **except** the 2-local cases.

Sophisticated computational methods of Holmes and Wilson allowed most maximal subgroups of \mathbb{M} to be determined by 2013.

mmgroup

Unfortunately, Holmes and Wilson's code is unreleased and still undesirably slow. Only in 2022 did fast group operations in \mathbb{M} become possible: Seysen's `mmgroup` Python package.

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except $2^{5+10+20} \cdot (S_3 \times L_5(2))$.

The Maximal Subgroups of \mathbb{M}

$2 \cdot \mathbb{B}$	$(D_{10} \times \text{HN}) : 2$	$(A_5 \times A_{12}) : 2$
$2^{1+24} \cdot \text{Co}_1$	$5^{1+6} : 2 \cdot J_2 : 4$	$(A_6 \times A_6 \times A_6) \cdot (2 \times S_4)$
$2^{2 \cdot 2} E_6(2) : S_3$	$(5^2 : 4 \cdot 2^2 \times U_3(5)) : S_3$	$(A_5 \times U_3(8) : 3) : 2$
$2^{2+11+22} \cdot (M_{24} \times S_3)$	$5^{2+2+4} : (S_3 \times \text{GL}_2(5))$	$(\text{PSL}_3(2) \times S_4(4) : 2) \cdot 2$
$2^{3+6+12+18} \cdot (\text{PSL}_3(2) \times 3 \cdot S_6)$	$5^{3+3} \cdot (2 \times \text{PSL}_3(5))$	$(\text{PSL}_2(11) \times M_{12}) : 2$
$2^{5+10+20} \cdot (S_3 \times \text{PSL}_5(2))$	$5^4 : (3 \times 2 \cdot \text{PSL}_2(25)) : 2$	$(A_7 \times (A_5 \times A_5) : 2^2) : 2$
$2^{10+16} \cdot \text{O}_{10}^+(2)$	$(7 : 3 \times \text{He}) : 2$	$M_{11} \times A_6 \cdot 2^2$
$3 \cdot \text{Fi}_{24}$	$7^{1+4} : (3 \times 2 \cdot S_7)$	$(S_5 \times S_5 \times S_5) : S_3$
$3^{1+12} \cdot 2 \cdot \text{Suz} : 2$	$(7^2 : (3 \times 2 \cdot A_4) \times \text{PSL}_2(7)) : 2$	$(\text{PSL}_2(11) \times \text{PSL}_2(11)) : 4$
$S_3 \times \text{Th}$	$7^{2+1+2} : \text{GL}_2(7)$	$U_3(4) : 4$
$(3^2 : 2 \times \text{O}_8^+(3)) \cdot S_4$	$7^2 : \text{SL}_2(7)$	$\text{PSL}_2(71)$
$3^{2+5+10} : (M_{11} \times 2 \cdot S_4)$	$11^2 : (5 \times 2 \cdot A_5)$	$\text{PSL}_2(59)$ 59:29
$3^{3+2+6+6} : (\text{PSL}_3(3) \times \text{SD}_{16})$	$(13 : 6 \times \text{PSL}_3(3)) \cdot 2$	$\text{PSL}_2(41)$
$3^8 \cdot \text{O}_8^-(3) \cdot 2$	$13^{1+2} : (3 \times 4 \cdot S_4)$	$\text{PGL}_2(29)$
	$13^2 : \text{SL}_2(13) : 4$	$\text{PGL}_2(19)$
	41:40	$\text{PGL}_2(13)$

The maximal subgroups of \mathbb{M} .

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The Choice of Tools

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Problem: permutation representations of 2-local subgroups seem elusive.

Could use `mmgroup` representations, but `mmgroup` currently has limited functionality.

Hybrid Groups

Hulpke (2024) handles the 2-local subgroup $2^{10+16} \cdot O_{10}^+(2)$ differently: as a *hybrid group* in GAP.

The implementation, based on Dietrich and Hulpke (2021), represents a group G by a solvable normal subgroup H , the quotient G/H , and the conjugation action of G/H on H .

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This works for $2^{5+10+20} \cdot (S_3 \times L_5(2))$ too: we can obtain the hybrid representation from `mmgroup`, and even translate between it and the `mmgroup` copy.

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Mostly straightforward: use GAP commands.

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Theorem (Brauer's Theorem on Induced Characters)

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Implementing this in GAP specifically for $2^{5+10+20} \cdot (S_3 \times L_5(2))$ is mostly OK, especially since elementary subgroups are so nice.

The Computation (II)

Problem: we need the character table of $P \in \text{Syl}_2(2^{5+10+20} \cdot (S_3 \times L_5(2)))$, which is huge: $|P| = 2^{46}$, and P 's character table is 26758×26758 .

Introduce some optimisations.

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Thank you!